

# Attacks against RSA signatures

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## 1 Fault attacks against RSA signatures

1. Implement the signature generation algorithm using the Chinese Remainder Theorem (CRT) using the Sage library. More precisely, to compute  $s = m^d \bmod N$ , compute

$$s_p = s \bmod p = m^{d \bmod p-1} \bmod p$$

and

$$s_q = s \bmod q = m^{d \bmod q-1} \bmod q$$

Recover  $s \bmod N$  from  $s_p$  and  $s_q$  using the CRT.

2. Assume that an error occurs during the computation of  $s_p$ , that is, an incorrect value  $s'_p \neq s_p$  is computed while  $s_q$  is correctly computed. Explain and implement how to recover the factorization of  $N$  from  $s$ , following the Bellcore attack [BDL97].
3. How could such error be detected ? Propose and implement a simple method to detect such error.

## 2 Optional: the Desmedt-Odlyzko attack

Implement the Desmedt-Odlyzko attack [DO85] described in the lecture, with the RSA signature scheme  $\sigma = H(m)^d \bmod N$ . The attack computes a forged signature as a multiplicative combination of existing signatures.

For simplicity the hash function can be computed as follows:

```
import hashlib

def sha1(s, digestsize=50):
    m = hashlib.sha1()
    m.update(s)
    return Integer(m.hexdigest(), base=16) % 2^digestsize
```

To detect smooth numbers among  $H(m_i)$ , one can use the `factor()` function from Sage. Given a  $(\ell + 1) \times \ell$  matrix  $M$  of exponent vectors modulo  $e$ , one can obtain a vector from the kernel of  $M$  using:

```
v=Matrix(GF(3),M).left_kernel().matrix()[0]
```

assuming that we work with public exponent  $e = 3$ . Using such vector  $v$  one can then express one row of  $M$  as a linear combination of the others. This enables to express one  $H(m_\tau)$  as a multiplicative combination of the others. Eventually, this enables to express one signature as a multiplicative combination of the others, hence a forgery.

To test the attack, one can use small parameters, for example `digestsize=50`, and a number of primes  $\ell = 100$ . It can be interesting to optimize the running time by varying  $\ell$  for a fixed `digestsize`. Eventually, one can experiment the attack for increasing values of `digestsize`.

## References

- [BDL97] Dan Boneh, Richard A. DeMillo, and Richard J. Lipton. On the importance of checking cryptographic protocols for faults (extended abstract). In *Advances in Cryptology - EUROCRYPT '97, International Conference on the Theory and Application of Cryptographic Techniques, Konstanz, Germany, May 11-15, 1997, Proceeding*, pages 37–51, 1997.
- [DO85] Yvo Desmedt and Andrew M. Odlyzko. A chosen text attack on the RSA cryptosystem and some discrete logarithm schemes. In *Advances in Cryptology - CRYPTO '85, Santa Barbara, California, USA, August 18-22, 1985, Proceedings*, pages 516–522, 1985.